### Admissibles and computations

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1. admissible ordinals?

Cantor a trouvé une loi d'engendrement de la multitude des nombres ordinaux finis et transfinis, il a trouvé une dynastie, celle des Aleph, et cela, à l'aide de deux principes seulement, l'un immanent (additif), l'autre transcendant (passage à la limite) : Cantor, législateur de l'infini. —Sinisealli, Horror vacui





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Equivalence class of well-orderings by isomorphism

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Transitive set well-ordered by  $\in$ 

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Ordinals + transfinite induction

measuring of provability strength

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Levels of constructibility  $L_{\alpha}$ 





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measuring of provability strength

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Levels of constructibility  $L_{\alpha}$ 

Coding infinite ordinals

 $\langle x,y \rangle = 1$  iff x < y ( $\langle \cdot, \cdot \rangle$  codage de  $\mathbb{N}^2$  dans  $\mathbb{N}$ ) ordre sur  $\omega$ 

ordre sur  $\omega$  *vs* ordre sur une partie de  $\omega$ (soit *n* est comparable à une infinité d'entiers, soit *n* est isolé) si  $\forall x \langle n, x \rangle = 0$ , alors soit *n* est le plus petit élément, soit il n'est pas dans l'ordre. si de plus  $\forall x \langle x, n \rangle = 0$  alors *x* n'est pas dans l'ordre

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Transitive set well-ordered by  $\in$ 

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An ordinal is recursive if it has a recursive coding on  $\omega$ 

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If  $\alpha$  has a recursive coding on  $E \subset \omega$  then it is recursive

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Coding infinite ordinals

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A code for  $\beta < \alpha$  extracted from a code for  $\alpha$ 









#### closed enough ordinals

#### limit, limit of limits, etc

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#### closed enough ordinals

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 $\omega_1^{\text{CK}}$  is the sup of the recursive ordinals

 $\begin{array}{l} \Sigma_0\text{-separation:}\\ \text{for any set } E, \ \Sigma_0 \ \text{formula } \varphi,\\ \text{there exists } X \subseteq E \ \text{such that}\\ x \in X \iff \varphi(x) \end{array}$ 

 $\begin{array}{c} \text{extensionality}\\ \text{induction}\\ \text{empty set}\\ \text{pairing}\\ \text{union}\\ \Sigma_0\text{-separation}\\ \Sigma_0\text{-collection} \end{array}$ 

#### closed enough ordinals

#### limit, limit of limits, etc

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 $\Sigma_0$ -collection:

for any  $\Sigma_0$  formula  $\phi(x, y)$  s.t.  $\forall x \exists y \phi(x, y)$ , we have that  $\forall X \exists E$ 

s.t.  $[e \in E \iff \exists x \in X \phi(x, e)]$ 

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 $\alpha$  is admissible if  $L_{\alpha}$  is a model of KP

admissible ordinals?

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 $\omega_1^{CK,r}$  is admissible for every real r

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2. gaps? *ITTMs, clockable ordinals, gaps.* 

#### Infinite time Turing machines (ITTM) Im sup special limit *state head* limit behaviour can compute on *reals*

ITTMs are extensions of Turing machines to transfinite time.

For simplicity: alphabet {0, 1} Ordinal stages

- successor ordinals: it works exactly as a Turing machine
- *limit* ordinals: each cell is set to the lim sup of its values, head is rewinded back to the origin and the machine enters a special limit state *L*
- = 3 tapes: input, scratch, output

A real (infinite binary string) can be considered as input (*oracle* computation) and output Example: a coding of an ordinal may be written, or taken as input by an ITTM

Infinite time Turing machines (ITTM) lim sup special limit state can compute on *reals* 

 $\alpha$  countable is *writable* if  $r_{\alpha}$  is



#### Infinite time Turing machines (ITTM)



Any recursive ordinal  $\alpha$  can be finitely represented by a Turing Machine  $\mu$  such that  $\mu(i) = r_{\alpha}(i)$ Among these TM we choose one (e.g. of smallest index)  $\alpha \mapsto \mu_{\alpha}$ Our order (partial) :  $\mu_{\alpha} \prec \mu_{\beta} \iff \alpha < \beta$ This order  $\prec$  is not recursive and is of type  $\omega_{1}^{CK}$ 

This order does not exhibit a recursive minimality - we will improve it using ITTMs



Infinite time Turing machines (ITTM) lim sup special limit state can compute on *reals* 

 $\alpha$  countable is *writable* if  $r_{\alpha}$  is

 $\alpha$  countable is *clockable* if  $\exists$  ITTM which halts exactly after  $\alpha$  many steps

head limit behaviour

Infinite time Turing machines (ITTM) Itm sup special limit state head limit behaviour can compute on *reals* 

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 $\lambda_{\infty} = \sup$  of writables  $\gamma_{\infty} = \sup$ 

 $\gamma_{\infty} = sup of clockables$ 

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 $\zeta_{\infty} = \sup$ 

 $\Sigma_{\infty} = \sup$ 

 $\lambda_{\infty} =$ sup of writables  $\gamma_{\infty} =$  sup of clockables

eventually writable ordinals accidentally writable ordinals

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Infinite time Turing machines (ITTM) Im sup special limit *state head* limit behaviour can compute on *reals* 

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There are non-clockable writable ordinals

 $\stackrel{\geq}{=}$  Admissibles are not clockable

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 $\stackrel{\geq}{=} \omega_1^{CK}$  starts a gap of length  $\omega$ 

 $\stackrel{\frown}{=}$  Admissibles are not clockable

 $\frac{1}{2} \omega_1^{CK}$  starts a gap of length  $\omega$ 

A gap length is a limit ordinal

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Is there a link between a gap's size and its starting point?

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How are gaps distributed?

Is there a link between a gap's size and its starting point?

Is there a gap with exactly one admissible ordinal properly inside?



3. a gap with exactly one admissible

## One admissible



### One admissible

 $L_{\lambda_{\infty}}\prec_{1}L_{\zeta_{\infty}}\prec_{2}L_{\Sigma_{\infty}}$ 

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There is a gap with at least one admissible inside

CL

There are many admissible ordinals between the sup of writables  $\lambda_{\infty}$  and  $\zeta_{\infty} < \Sigma_{\infty}$ .

 $\exists \alpha < \beta < \gamma \text{ s.t. } \beta \text{ adm. and } L_{\gamma} \text{ witnesses } [\alpha, \beta) \text{ contains no clockable}$ 

This is a  $\Sigma_1$  statement, verified in  $L_{\Sigma_{\infty}}$ , and thus also in  $L_{\lambda_{\infty}}$ . Let  $\alpha, \beta, \gamma < \lambda_{\infty}$  witness this.  $L_{\gamma}$  for 0, believes that  $\beta$  is an admissible ordinal properly inside a gap.

Since any ITTM-computation on the empty input of length  $< \beta$  is already contained in  $L_{\beta}$ ,  $[\alpha, \beta)$  is indeed inside a gap which properly contains the admissible ordinal  $\beta$ .

NCL

Design an algorithm that checks if a real x is a code for an ordinal which is the starting point of a gap containing  $\omega_1^{CK,x}$ .

This algorithm accepts any code for  $\lambda_{\infty}$ .

There is an x that is accepted by this algorithm

### is a $\Sigma_1$ statement.

By the  $\lambda - \zeta - \Sigma$  theorem, we have a witness of this property in  $L_{\lambda_{\infty}}$ , which has to be the code of an admissible ordinal  $< \lambda_{\infty}$  beginning a gap with an admissible properly inside.

### One admissible

 $L_{\lambda_{\infty}} \prec_1 L_{\zeta_{\infty}} \prec_2 L_{\Sigma_{\infty}}$ 

and a second

There is a gap with at least one admissible inside

The first gap with at least an admissible inside will only have one admissible  $\tau$ , and ends at  $\tau + \omega$ 

### One admissible

#### Algorithm :

*Primary gap detection* : we run a universal online-simulation of all ITTM in order to detect gaps. We thus observe gaps  $\omega$  steps after their starting point.

At each starting points  $\alpha$  we have at our disposal a coding of  $\alpha$ .

We then start a secondary gap detection with oracle  $\alpha$  until either :

- the primary gap ends; we then stop secondary detection we continue the primary detection until next gap
- the first secondary gap is observed; we then halt.

#### Proof :

here is a gap with at least one admissible inside

If no admissible inside a gap, then it halts after  $\lambda_{\infty}$ . Contradiction. Halting time:  $\alpha + \omega_1^{CK, \alpha} + \omega = \omega_1^{CK, \alpha} + \omega$ . Hence this is the end of the gap. 4. beyond one admissible

The first gap with at least  $n < \omega$  admissibles inside will only have *n* admissibles inside, and ends  $\omega$  steps after the last one

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Direct adaptation of previous proof

Gap:  $[\alpha_0, \ldots, \alpha_1, \ldots, \ldots, \alpha_n]$ 

Computation of  $\omega_1^{CK}$  relativized to detect the next admissible :

 $\alpha_{i+1} = \omega_1^{CK, \alpha_i}$ 

The first gap with at least  $n < \omega$  admissibles inside will only have *n* admissibles inside, and ends  $\omega$  steps after the last one

 $(\alpha < \lambda_{\infty})$  After the writing time for  $\alpha$ , the first gap with at least  $\alpha$  admissibles inside has exactly  $\alpha$  inside

The first gap with at least  $n < \omega$  admissibles inside will only have *n* admissibles inside, and ends  $\omega$  steps after the last one

Not a direct adaptation of previous proof because of limits of admissibles

Problem: if  $\alpha_0, \alpha_1, \ldots, \alpha_n, \ldots \rightarrow \alpha_{\omega}$  then sometimes  $\alpha_{\omega}$  is admissible (*recursively inaccessible*) and sometimes not.



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Ranksrank-0: admissiblesrank- $\alpha$  + 1: admissible limits of rank  $\alpha$  adm.rank- $\lambda$ : adm. limits of rank  $\beta$  adm.  $\forall \beta < \lambda$ 

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- $_{\geq}$  Let  $f: \omega_1 \to \omega_1$  be ITTM-computable. Then there is an
- = ordinal  $\alpha$  starting a gap of size  $\geq f(\alpha)$

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... and a few other improvements

5. sacks' characterizations of admissible ordinals

Il y a en nous une sensation finie de « l'infini». Et ce n'est qu'un effet – une conséquence. Ce n'est pas une preuve de quoi que ce soit. —Paul Valéry, 1910





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For every countable admissible  $\alpha$ ,  $\exists r$  s.t.  $\alpha = \omega_1^{CK,r}$ 

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How can we understand & prove this in an explicit way?

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simulation of all ITTM's w-online  $\omega$  machines halt for a given simulation time, we choose the first one in the simulation  $\mu \prec \nu \iff (\mu \text{ and } \nu \text{ are chosen and } \mu \text{ halts before } \nu)$ if we run the above process up to  $\alpha$  we get an order for all clockables before  $\alpha$ first run of  $\omega$  non clockables  $\Longrightarrow \omega_1^{CK}$ the order type of clockables below a gap is exactly the starting point of the gap coding of  $\lambda_{\infty}$ same with oracle (or input) coding of  $\omega_1^{CK,r}$ in some sense these are the simplest codings

For every countable admissible  $\alpha$ ,  $\exists r$  s.t.  $\alpha = \omega_1^{CK,r}$ 

How can we understand & prove this in an explicit way?

The writing time of any  $\alpha < \lambda_\infty$  is the sup of all ends of gaps that start before  $\alpha$ 

6. building up to a proof

*Qui est là ? Ah très bien : faites entrer l'infini.* 

—Aragon, Une vague de rêves, 1924

### Successor admissible case

6

Successor admissible case  $\alpha = \beta^+$  ( $\alpha$  and  $\beta$  are admissible) Code  $\beta$  in a real *r* to obtain  $\omega_1^{CK,r} = \alpha$ 



#### Successor admissible case

If  $r_{\beta}$  codes an ordinal  $\beta$ ,  $\omega_{1}^{CK,r_{\beta}} \geqslant \beta^{+}$ 

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We have to make sure that r does not code more...

Find the lowest simplest code
### Successor admissible case

If coding level of  $r_{eta}$  is  $<eta^+$ , we have equality and everything works.

Otherwise, we need to do something a little bit more sophisticated.

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Recursively inaccessible case

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Consider the first recursively inaccessible  $\iota_0$ 



We need some tool to construct r

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 $\{r_i : i \in \omega\} \text{ s.t. } \forall i, r_i <_{\mathcal{T}} r_{i+1} \\ \exists r \text{ s.t. } \forall i, r_i \leqslant_{\mathcal{T}} r \text{ but } \bigoplus_i r_i \nleq_{\mathcal{T}} r \leqslant_{\mathcal{T}} (\bigoplus_i r_i)'$ 

*r* needs to code the  $r_i$ 's but in such a way that  $\forall i, r_i$  can be computed from *r* but not uniformly.

*r* is a *mapping* from  $\omega^2$  to {0, 1} that contains  $r_i$  in the  $\mathfrak{b}(i)$  column.

We ensure that  $\bigoplus_i r_i \not\leq_T r$  by adding the needed information to hide the  $r_i$ 's.

*r* is constructed as  $\bigcup_i o_i$  where the  $o_i$ 's are compatible oracles that represent the left part of *r* up to column  $\mathfrak{b}(i)$ 

We assume that  $o_i$  has been built and that we know b(i), and we give the construction for  $o_{i+1}$  and b(i+1).

Consider all  $\varphi_i^{\tau}(\langle i+1,j \rangle)$  computations for every *j* and every finite extension  $\tau$  of  $o_i$ .

Look for the first  $(\tau, j)$  such that it either (i) converges and is  $\neq r_{i+1}(j)$ , or (ii) for every extension of  $\tau$ , it diverges.

b(i + 1) is then taken to be greater than [case (*i*)] the maximum between b(i) and the greatest column reached during the computation, [case (*ii*)] the greatest column reached in the enumeration of the extensions of  $o_i$ ; which is the column from which every extension will make the computation diverge on *j*.

We need some tool to construct r

 $\{r_i : i \in \omega\} \text{ s.t. } \forall i, r_i <_{\mathcal{T}} r_{i+1} \\ \exists r \text{ s.t. } \forall i, r_i \leqslant_{\mathcal{T}} r \text{ but } \bigoplus_i r_i \notin_{\mathcal{T}} r \leqslant_{\mathcal{T}} (\bigoplus_i r_i)'$ 

The *lost* lemma is actually very close to an old classical computability result

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ACADÉMIE DES SCIENCES.

8° Republica Argentina. Publicaciones de la Comision nacional de la energia alo no a disce a companya de la companya de la companya de la energia

Il signale également un fascicule polycopié : Contribution du Laboratoire d'astronomie de Lille, nº 2. Numéro spécial, à l'occasion du Colloque international de Liége sur Les particules solides dans les objets astronomiques.

ARITHMÉTIQUE. — Sur le semi-réseau constitué par les degrés d'indécidabilité récursive. Note (\*) de M. Daniel Lacombe, présentée par M. Émile Borel.

Extension de certains résultats de Kleene et Post (1) et solution de quelques questions posées par ces auteurs (3).

Nous utiliserons dans ce qui suit les définitions et les notations de S. C. Kleene-E. L. Post (<sup>1</sup>). Nous désignerons par D l'ensemble des degrés d'indécidabilité (ou, pour abréger : degrés) et par D, l'ensemble des degrés arithmétiques ( $D_A \subset D$ ). D et  $D_A$  sont munis d'une relation d'ordre partiel (notée  $\leq$  et  $\leq$ ) et d'une opération  $a \rightarrow a'$  partout définie (nous désignerons par  $a^{ii}$  le résultat de cette opération ifèré i fois à partir de a).

Les deux théorêmes suivants se démontrent au moyen des méthodes classiques (fondées essentiellement sur la *forme normale de Kleene*) complétées par l'utilisation de fonctions majorantes. Les conditions (A) et (B) du théorême l constituent deux cas particuliers d'une condition plus générale que nous ne pouvons énoacre ic, et qu'il serait d'ailleurs intéressant d'élargir.

THEOREME I. — Soit  $S = u_s, u_1, \ldots, u_n, \ldots$  une suite infinie de degrés, strictement croissante (c'est-à-dire telle que i < j entraîne  $u_i < u_j$ ) et satisfaisant à l'une ou l'autre des conditions suivantes :

(A) il existe un degré a tel que, pour tout i,  $u_i = a^{(i)}$ ;

(B) il existe un degré b tel que, pour tout i,  $b < u_i < b'$ .

Soit U l'ensemble de degrés défini par la condition :



<sup>(\*)</sup> Séance du 27 octobre 1954.

(\*) Ces questions sont en général indiquées, dans l'article cité ci-dessus, par le signe : 3, UD lo lequel renvoie à une Note de la page 380. Les résultats annoncés dans cette Note n'ont pas encore, à notre connaissance, cét publiés. SÉANCE DU 3 NOVEMBRE 1954.

(1) pour tout degré x, on a

 $(x \leq d_1 \ et \ x \leq d_2) \Leftrightarrow x \in \mathbb{U};$ 

(2) c ne vérifie aucune des deux inégalités

 $c \leq d_1$  ct  $c \leq d_2$ .

Remarque 1. — Ce théorème montre que S (ou, ce qui revient au même, U) ne possède pas de borne supérieure précise.

Remarque 2. — La relation (1) montre que le couple  $(d_i, d_i)$  ne possède pas de borne inférieure précise. De l'existence de suites S satisfaisant à (A) ou (B) on déduit donc immédiatement que D ne constitue pas un réseau [résultat démontré par Kleene et Post au moyen d'une suite de type (A)].

Remargue 3. — Kleene et Post ont montré l'existence de suites 5 satisfaisant à la condition (B), le degré b appartient à D<sub>4</sub>, il en est de même pour b' et pour tous les  $u_i$  (et l'on a U  $\subseteq$  D<sub>1</sub>). Cela n'entraine pas forcément que d<sub>i</sub> et d<sub>4</sub> puissent ètre pris eux, cet l'on a U  $\subseteq$  D<sub>1</sub>). Cela n'entraine pas forcément que d<sub>i</sub> et d<sub>4</sub> puissent être pris eux, causi dans D<sub>4</sub>. La relation (1) montre en effet que l'ensemble U est entièrement déterminé par la donnée de d<sub>i</sub> et d<sub>5</sub>. Or D<sub>4</sub> est dénombrable (donc aussi l'ensemble des couples formés de deurés arithmétiques). Mais Kleene et Post ont montré que les ensembles tels que U, déterminés dans D<sub>4</sub> par des suites S de D<sub>4</sub> croissantes et auisfaisant à (B) (avec b dans D<sub>4</sub>), forment une famille ayant la puissance du continu. Il en résulte que pour certaines de ces auites il n'existe aucun couple (d<sub>1</sub>, d<sub>2</sub>) formé de degrés arithmétiques et saitafissant à la relation (1). Le théorème suivant donne une condition suffisante pour l'existence d'un tel couple.

Étant donnée une suite de degrés quelconque  $S = u_i, u_1, \ldots, u_n, \ldots$  et une fonction  $\varphi$  de deux variables (entières  $\geq o_i$ ), nous dirons que  $\varphi$  étumére S si, pour tout i, la fonction d'une variable  $\varphi_i$  définie par  $\varphi_i(x) = \varphi(i, x)$  est de degré  $u_i$ .

THEOREME II. — Si, dans le théorème I-hypothèse (B), le degré b est arithmétique et si la suite S peut être énumérée par une fonction arithmétique, alors les degrés d, et d, satisfaisant aux relations (1) et (2) peuvent être pris (d'une infinité



Remarque 3. — Les méthodes de Kleene-Post permettent de déterminer des suites S satisfaisant aux conditions de ce théorème II. Il en résulte que  $\mathbf{D}_A$  ne constitue pas un réseau.

<sup>(1)</sup> Ann. Math., 59, 1954, p. 379-407.

1108

ACADÉMIE DES SCIENCES.

S° Republic Argentina. Publicaciones de la Comision nacional de la energia

Il signale également un fascicule polycopié : Contribution du Laboratoire

d'astronomie de Lille, nº tional de Liége sur Les

> ARITHMÉTIQUE. dabilité récursive M. Émile Borel.

Extension de certain questions posées par ces

Nous utiliserons da S. C. Kleene-E. L. Poi d'indecidabilité (ou, p arithmétiques ( $\mathbf{D}_{\mathbf{L}}, \mathbf{D}_{\mathbf{L}}$ ) (notéo  $\leq t \leq 2$ ) et d'u par d'u la résultat de co Les deux théorèmes siques (fondées essenti par l'utilisation de lo nous ne pouvous énonc

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pace de suites S satisfaisant egré b. Lorsque, dans cette at de même pour b' et pour 'cément que d, et d, puissent s en effet que l'ensemble U  $d_{\gamma}$ . Or D, est dénombrable degrés arithmétiques). Mais B) (avec b dans D<sub>A</sub>) forment suite que pour certains de de degrés arithmétiques et donne une condition suff-

 $= u_0, u_1, \dots, u_n, \dots \text{ et une}$ s dirons que  $\varphi$  énumère S si, par  $\varphi_i(x) = \varphi(i, x)$  est de

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### ANNALS OF MATHEMATICS Vol. 64, No. 3, November, 1956 Printed in U.S.A.

### **ON DEGREES OF RECURSIVE UNSOLVABILITY\***

BY CLIFFORD SPECTOR

(Received August 8, 1955)

In Kleene-Post  $[4]^1$  a number of questions concerning the structure of the upper semi-lattice of degrees were left unanswered. The present paper contains the answers to those questions under the scope of [4] Footnote 3. With the exception of the density problem ([4], 2.2), the methods used are variations of those developed in [4]. The construction employed in showing that the degrees are not dense involves a general ratio of the methods of [4], and constitutes the main result of this paper (Theorem 4) that there are minimal degrees of recursive unsolvability.<sup>2</sup> Familiarity with [4] is assumed.<sup>3</sup>

of Bells." TOTICE OF CHE MERGEROL

He also made concert arrustees Amer- rangements of such songs as DONOGHUE-UNTERWEISER-Mr and Mrs "The Old Chisholm Trail." "Old Ex-Paint" and "Whoopee-ti-vi-vo. shville Git Along, Little Dozies." In 1925 he put the folk song. 'Home on the Range," into Coun-

sheet-music form.

intain CLIFFORD SPECTOR, of The **MATHEMATICIAN. 30** 

to his and a Special to The New Yort Times. Evans

PRINCETON, N. J., July 29-Dr. Clifford Spector, Associate Professor of Mathematics at the

University of Michigan, died to-29:day of a cerebral hemorrhage at

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United States mathematicians to attend an international conference on symbolic fogic at the Warsaw Acadeby of Science. A KER Last year he did ma remarical research at the Institute for Allvanced Studies here, 29 He leaves his wife, the former vice Lea Eisner; two children. Alan th & and Judith: his parents, Mr. and Mrs. Robert Spector of New BADINELLI-Juna, at 50 York, and a brother. Gilbert, New Professor of Music at Kansas

to Lester Lewis Borden, son of Mr. and -----Mrs. George Bloom of Great Neck. The wedding was held at the Hatel Pierre, on

Louis Gross announce the marriage of their laughter, Mrs. Jeache G. Unterweiser, to Edward A. Donoghue, on July 19, in HUJSA-KATZ-Mr. and Mrs. Adolph Katz of

Woodmere, L. L. ioyfully announce the St. and Amsterdam Ave. marriage of their daughter, Carole Joyce, GOLDFEDER-Simon, beloved husband to Charles Jay Huisa. Ceremony took place at Temple Israel, Lawrence, L. L. on Sun.

ECHI DELL'ARIA Dr. and Mrs. Saluatore Dell'Aria of 103 E. \$6th St. announce the marriage of their daughter, Marina, to Charles J Segui, at the St. Ignatius Loyola Thurs., July 27, 1961 BLAPIRO-BCHAPIRO-Mr. and Mrs. Morris Schapiro, Lake Success, N. Y., announce GRIFFIN-Elizabeth Willett Ryle, on July 28

the marriage of their daughter, Marcial Eileen, to Dr. Herman Shearo, son of Mr. and Mrs. Gershon Sheairo, Baysice,

### Anniversaries

### Apaths

hoe and Alice Stunmer, Brdher or Lonal Koight and Norman 1. Allen, also survived by ten grandchildren, service at the Fairchild Chapel, Franklin Ave, at 12th HARF-Emilie, beloved wile of the idia nax, nations to your Church Charity would be

found sourow the death of its trustee, HATRY-William A., beloved huband of the Chester A. Ailen.

RAYALIND H EIERC, Presiden -Margaret Sansoney, & Greenwich 1961, wife of the late Allen and sister of Mrs. Joseph A. Lee and Mrs. Howard W. Clark, Serv-1 icas at First Presbyterian Church, Green-1 wich, on Monday, at 11 A. M. Interment St. Louis, Mo. Please omit flowers. ANUERSON-James E., belover husbaad of

actime Kennedy Angelson of 30 Park HOFFMAN-Rose, devoted mother of Ray East, Euneral from the William Fromm, Joan Behrman and Michael, Healy & Son Funeral Home, 271 War Yonkers, on Morday, at 0:20 ur Lady of the Rosary, Yonkers, at 10 ivel, beloved husbard of Estre

Y., beloved with of the latel Ave Reposing at the Westchester Main SL, recease, Requestor al Maine Westchester al Maine SL, recease, Requirem Mass in the

Brudner, Services Monday, 9:30 A. M. Diverside " 76th St. and Amsterdam Ave

FRIEND-Cecilia beloved mother of Priscilla O'BRIEN-EIIZauein Charlotte, Lea-Fan Rosen and heloved sister of Albert and Louis Voloin and Shirley Kirschenhaum, and loving grandmother, suddenly, on Fire Is and, Sat. July 29, Services Mon. July 31, at 1 P. M. at "The Riverside." 76th

Leah, devided father of Elinabeth, Karin OFFENTHER-Ivider, beloved husband o Dovler, Arthur, Eric and Marc: brother rannie, devided father of Harola; dariins of Balle Hoke and Irving, Service Sunrday, 1 P. M., at the Park West, 70th and Columbus GRAHAM-Mora O'Brien

Whidden Graham, mother of W. Crost Graham, Funeral private, Interment Balt more Ald

widow of Henry Artnur Grinne, dauguler of I grews Hospital Fund, Booln Bay Harbor the rate William and Mary Elizabeth Ryle, PALMER-Minnie Y., beloved mother of Mrs. Allred Elv. William and Albert Biuton Strange, Funeral service at the Chapel of St. James Church, Maoison Ave. at /ist St., on Mon., July 31, at 12 

KICHOLDSUH Ave., Bx. Funera Hum Moss.

NO. FT. Harrison, Clearwaler, Fla. GLYNN-Happy Silver Anniversary to my HARDT-Marie Do, (ree Dinketsiel), on July dearest earchts. Esther and too, Win love 29, 1931, of 47 Wooobine Ave, Larchmont, always, your own, HONEY ANN. Y, beloved wite et John P., loving mother of Erminy Marks and Louise Bauman. grandmother of Paul. Reposing at the FoxIPLUINIKOFF-Jack, beloved husband Funeral Nome, 60 Post Rd. Larchmont, Sophie, devotes latner of Davi and Su until Monday, 10 A. M. Religious service at Platt devites rather and fear brot In lieu of Nowers kindly send donations to practice PLL, John L. We somewfully an-

Cancer Fund, HARDI-Marie D. The Larchmont Temule Marie Hardt, and extends synpathy to her

Mrs. MAURICE MERMEY, President, dear mother of John F. Muller and Lise Alexancer, loving grandmotive and craft RATHEAACHER-Andree, of 54 Ocean grandmototier. Services Alondar, i P. M., vard, Atlantic Highlands, on Jery 22 "The Riverside", 76th 51, and Amster-wite of August Rathemacher, mot

tate May, devoted Lather of Patricia and Harry, deal buoine or harry A, and Grace Israel, Noice of funeral later (51), beloved Hucker-Jonn J. on July 27, 196, beloved husband of Mary (nee O'Hara), Funeral from Walter B. Cooke Funeral Home, 2135 Westchester Ake, Monday, 7:30 A. M.

Solemo Requiem Mass St. Raymond's A. M. Interment Calvery

ing grandmother, Services Sun, 10 A. Mass of Requirem at the Church of HOWARD—Luella, on July 28, 1961, in her ady of the Rosary, Yinkers, at 10, 93rd year, formerly of 2281 Loring PL. A manual sector of the sector

dear tainer of Meyer, Corplia, Rebecca and JACKSUN-Harry, beloved husbani of Rebecca, nancis; pergecu grandiather of Edward, devoted father of Philip, dear brother, routing grandtather. Services Suncay, 12:45 P. M., "The Riverside." 76th St and Amsterdam

or an industry of Irona Fleutiei ann JACKSON-Harry. The Bronx Chapter of In ties of flowers, please contribute to the Hadassah records with deep sorrow thei Damin Runyon Fund Haddassan records with deep solution inc. Dates the barrier to the solution of Mary, passing of Harry Jackson, belowed husband RUMMEL—Harold, lowing husband of Mary, of Buddy Jackson, our estremed Vice davide fathe of Louise, daving bother

Auna and Martin P. AL, "Ine Riverside, ioth St. and Amstercam Ave

J., formery of Newark. wile of join 1. O brich and mother of Euzabeth M. and Rev. John E. O'brich. Funeral from the Joines J. Higgins and Son Murtuars, 414 Westminister Acc., Enza bern, on Lues., Aug. 1, ar 9.30 A. M. Solemn Requirem mass Immaculate Con-ception Crierco, 10 A. M.

rangiaber. Services today, 10.30 A.M. "heiman's Euneral Chicels," last Grand Concourse, Bronx

OLCO (-Alfeed Van Santynard, July 78 Sur vived by his wile, fwo sons, a daughter and sty grandenilgren, Services Motoay, July 31 A. M., Dicolt residence, West Southoot Maine. Donelions may be male to st.

Arthur M. Loving, granamatier of Leslie Finke, auored sister of Certruse R. Asto-Seima Lasin and Samuel S. Shapiro arc devoied auni, services Mondey, 2 P. "Rivers de Memorial," 310 Coney II Avr., Hrookisn

PESNER-Minnie, beloved mother of Joseph Lily Katz and Herman, uear grandmomen A. M., Garlick's 'Parkside.' 1700 cure Island Ave. (Ave. N), Brooklin,

anie, devoted rather of David and Sidney Interreg Fri., July 28, 1961.

RASHKER—Jack L. We sorrowfully a nounce the bassing of our menber. West side institutional synglogue. LESTER UDELL, President.

PADA-Pudi, nationally known chotographer.

beloved nuscand of Annette and loving son Jeanette Rada, died sudderly, July 23. at South Miami, Fla

wife of August Rathemacher, mother of John and Richard, daughter of Mrs. Andree otiered for the recose of her soul at Friends may call at her residence Sunday atternoon and evening. Internent Mount Olivet Cemetery, Middletown Toyosnip, N. J. RICKLIN-Dr. Abraham, beloved husband of Mary, devoted taluet of Lytiling hermening Fasaler, cherished grandfather of brother of Dr. Joseph, Henry, and Haman Services at "The Bouleward," 312 Concy. Island Ave. (at Prospect Park), Browling

and the stall of the Coney Islung Hosnita Lydney, dear sister of Charlette Austo Alvin and Warren Newman, and the late Adelaide Lever. Services at "Fark West," 15 W. 79th St., Sun., July 30, 2:30 P. M.

darling sister of Cella Eliastern, Harry \_\_\_\_\_ defining selection of clina states where the selection of Park Circle, Brootling

### Card of Chanks.

and stacious exercisions of sympathy from the relations

The Chapel Of The Four Chaplains, in the spirit of the selfless men who act as its inspiration. is dedicated to serving people of all faiths. Every way in which we serve is aided by the most modern methods available; combined with the old-fashioned virtues of courtesy and dignity.

When you choose Universal, you select men and women who are sincerely interested in you and who are prepared to give you their very best efforts. This is what you have yuk ret al Universal.

EUNERAL CHAREL OF THE

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Schapiro, Lake Success, N.Y., announce GRIFFIN-Elizabeth Willett Rviz, on July 28

### Anniversaries

HARDI-Marie D. The Larchmont Temple

a and b form an exact pair for a degree set C if

both are above all degrees in C

univers

any degree below both is also below some degree in  ${\cal C}$ 

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### Anniversaries

# **a** and **b** form an exact pair for a degree set C if

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Every countable set of degrees in which every pair of elements is bounded has an exact pair

### $ι_0$ : *ω*-sequence $β_i$ of admissibles

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 $\beta_i$  codable by  $r_i$  in  $L_{\beta_i+1}$ 

6.-building up to a proof

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Where are the codes for the sequence members?

6. building up to a proof







There are arbitrarily long gaps in *L* where no news reals appear *Putnam* gaps

Let  $\beta > \alpha$  be countable ordinals such that there is an elementary embedding  $j: L_{\beta} \to L_{\omega_2}$  with *critical point cr*(j)  $\geq \alpha$ .

For every  $\gamma < cr(j)$ ,

 $L_{\omega_2} \models$  "No new reals appear between ranks  $\omega_1$  and  $\omega_1 + \gamma$ ."

No new reals thus appear between cr(j) and  $cr(j) + \gamma$ , by elementarity and absoluteness.



There are arbitrarily long gaps in *L* where no news reals appear *Putnam* gaps

If new reals appear at  $\alpha + 1$ , then among them is an arithmetical copy  $E_{\alpha}$  of  $L_{\alpha}$ 

 $E_{\alpha}$  is an arithmetical copy of  $L_{\alpha}$  if there is one-one function *f* from  $L_{\alpha}$ to  $\omega$  (and onto the field of  $E_{\alpha}$ ) such that  $\forall x, y \in L_{\alpha}$ ,  $x \in y \iff \langle f(x), f(y) \rangle \in E_{\alpha}$ 



If new reals appear at  $\alpha + 1$ , then among them is an arithmetical copy  $E_{\alpha}$  of  $L_{\alpha}$ 

Later seen as Jensen's mastercodes

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*r* is a master code for  $\mathcal{J}_{\xi}$  if  $\{x \subseteq \omega : x \leq_{\mathcal{T}} r\} = \mathcal{J}_{\xi+1} \cap \mathcal{P}(\omega).$   $E_{\alpha}$  is an arithmetical copy of  $L_{\alpha}$  if there is one-one function *f* from  $L_{\alpha}$ to  $\omega$  (and onto the field of  $E_{\alpha}$ ) such that  $\forall x, y \in L_{\alpha}$ ,  $x \in y \iff \langle f(x), f(y) \rangle \in E_{\alpha}$ 

 $\begin{array}{l} \mathcal{J}_0 = \varnothing, \ \mathcal{J}_{\omega\cdot\xi} = L_{\xi}, \ \mathcal{J}_{\omega\cdot\xi+n} = \\ \Delta_n(L_{\xi}), \ \text{where } \lambda \ \text{is a limit ordinal} \\ \text{and } \Delta_n(X) \ \text{is the set of all subsets} \\ \text{of } X \ \text{definable with parameters in} \\ \langle X, \in \rangle \ \text{by both } \Sigma_n \ \text{and } \Pi_n \ \text{first order formulae.} \end{array}$ 

There are arbitrarily long gaps in *L* where no news reals appear *Putnam* gaps

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New reals appear at level  $\xi$  for  $\mathcal{J}$  iff there is a master code for  $\xi$ . Furthermore, if *r* is a master code for  $\xi$ , then *r'* is the master code for  $\xi + 1$ .

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There are arbitrarily long gaps in *L* where no news reals appear *Putnam* gaps

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There exists  $\alpha$  such that  $L_{\alpha} \prec L_{\omega_1}$ ,  $\alpha$  is thus not definable in  $L_{\omega_1}$ . There is a countable  $v > \alpha$  such that  $L_v \prec L_{\omega_1}$ , and  $\alpha$  is already not definable in  $L_v$ .

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We can characterize the least such definability gap

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 $\begin{aligned} & E_{\alpha} \text{ is an arithmetical copy of } L_{\alpha} \text{ if} \\ & \text{there is one-one function } f \text{ from } L_{\alpha} \\ & \text{to } \omega \text{ (and onto the field of } E_{\alpha} \text{) such } \\ & \text{to } \omega \text{ (and onto the field of } E_{\alpha} \text{) such } \\ & \text{that } \forall x, y \in L_{\alpha}, \\ & x \in y \iff \langle f(x), f(y) \rangle \in E_{\alpha} \end{aligned}$ 

The ordinal  $v_0$ , which is the least v such that there is an ordinal  $\alpha$  not definable in  $L_v$ , can be characterized as the least  $\eta$  such that there exists an ordinal  $\delta < \eta$  such that  $L_\delta \prec L_\eta$ .



New reals appear at level  $\xi$  for  $\mathcal{J}$  iff there is a master code for  $\xi$ . Furthermore, if *r* is a master code for  $\xi$ , then r' is the master code for  $\xi + 1$ .

There are countable ordinals  $\alpha, \upsilon$  such that  $\alpha$  is not de-

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We can characterize the least such definability gap

 $\langle X, \in \rangle$  by both  $\Sigma_n$  and  $\Pi_n$  first order formulae.

 $\upsilon_0$  is clearly  $\leqslant$  the least such  $\eta$ ,  $\eta_0$ , since whenever one has  $L_{\alpha} \prec L_{\beta}$ ,  $\alpha$  is not definable in  $L_{\beta}$ .

Now, suppose that  $v_0 < \eta_0$ , in other words, for all  $\delta < v_0$ ,  $L_{\delta} \not\prec L_{v_0}$ . By Löwenheim-Skolem there is a countable elementary submodel of  $L_{110}$ . Take the  $\subseteq$ -least such model M. By the Condensation Lemma, there is an  $\alpha < v_0$  and an isomorphism j such that the Mostowski collapse of M is isomorphic to  $L_{\alpha}$  via *j*. *j* cannot be trivial as this would mean that  $L_{\alpha} \prec L_{\delta}$ , although  $\delta < v_0$  and  $v_0$ is the least such ordinal. We can thus consider K, the critical point of *j*. Since  $L_{\alpha} \cong M \prec L_{v_0}, L_{\kappa} \prec L_{i(\kappa)}$ . But then  $\kappa$  cannot be definable in  $L_{i(\kappa)}$ , and thus  $\upsilon_0 \leq i(\kappa)$ . But  $i(\kappa) < \upsilon_0$ , contradiction.

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We can characterize the least such definability gap
$\alpha$  definable at  $\gamma$  if definable without parameters in  $L_{\gamma}$  $\alpha$  codable at  $\gamma$  if appears in  $L_{\gamma+1}$  a real coding  $\alpha$  $\alpha$  countable at  $\gamma$  if  $L_{\gamma} \models "\alpha$  is countable".

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Consider  $\kappa = \aleph_{\alpha}$ .  $\kappa$  is definable as the greatest cardinal in  $L_{\kappa^+}$ . (Here  $\kappa^+$  denotes the least ordinal of cardinality greater than  $\kappa$ .) And thus  $\alpha$  is also definable in  $L_{\kappa^+}$ .

Löwenheim-Skolem's theorem, in conjunction with Mostowski's lemma and the Condensation Lemma, provides the countable  $\beta$  such that  $\alpha$  is definable in  $L_{\beta}$ .

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There is an upper bound for the ordinals that remain definable from some point on

 $\alpha$  definable at  $\gamma$  if definable without parameters in  $L_{\gamma}$ 

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*Memorable ordinals*: ordinals  $\alpha$  for which there exists  $\beta$  such that for any countable  $\gamma \ge \beta$ ,  $\alpha$  is still definable at  $\gamma$ .

Any countable  $\tau$  such that  $L_{\tau} \prec L_{\omega_1}$  is such an upper bound: if  $\alpha$  is definable at  $\beta$ , take  $\delta$  above  $\tau$  and  $\beta$  such that  $L_{\delta} \prec L_{\omega_1}$ . We then have  $L_{\tau} \prec L_{\delta} \prec L_{\omega_1}$ .  $\alpha$  is thus definable at  $\delta$ , since  $\delta$  is above  $\beta$ , and also at  $\tau$ .  $\tau$  is therefore above  $\alpha$  and any other definable ordinal. In fact, the least non-memorable ordinal  $\tau_0$  is the least ordinal  $\tau$  with uncountably many elementary extensions  $L_{\tau} \prec L_{\gamma}$ .

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# Codability

### Reals coding countable ordinals

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building up to a proof

## Codability

#### Reals coding countable ordinals

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Codes appear in an increasing way, sometimes in chunks

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## Codability

### Reals coding countable ordinals

Codes appear in an increasing way, sometimes in chunks

If  $\alpha$  and  $\beta$  are in the same codability gap and  $\alpha \ll \beta$ then  $r_{\alpha} <_{\tau} r_{\beta}$ 

#### Successor admissible case

Sime in

#### Successor admissible case

- marine

$$\omega_1^{CK,r_\beta} \ge \alpha$$
 but  $r_\alpha \not\leq_T r_\beta$ , so  $\omega_1^{CK,r_\beta} = \alpha$ 

 $21/\omega$ 

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Recursively inaccessible case

Everything happens in  $L_{\gamma}$ 

where the first code for  $\alpha$  appears

7. sacks' and jensen's theorems revisited

Est-il possible de raisonner sur des objets qui ne peuvent être définis en un nombre fini de mots ? Est-il possible même d'en parler en sachant de quoi l'on parle, et en prononçant autre chose que des paroles vides ? Ou au contraire doit-on les regarder comme impensables ? Quant à moi je n'hésite pas à répondre que ce sont de purs néants.

-Poincaré, La logique de l'infini, 1909

# How far does it work?

For every admissible  $\alpha < \omega_1^L$ ,

there exists a real r s.t.  $\alpha = \omega_1^{CK,r}$ 

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Optimality results concerning the degree of r





Sacks' theorem but for a sequence of admissibles?

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sequence of  $\alpha$  admissibles compatibility with the sequence of the first admissibles?

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Compatibility hypothesis: having each member admissible relative to the initial segment

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Compatibility hypothesis: having each member admissible relative to the initial segment

 $\begin{array}{l} \langle \alpha_{\beta} : \beta < \gamma < \lambda_{\infty} \rangle \text{ sequence of admissibles } < \omega_{1}^{L} \text{ s.t.} \\ \forall \delta < \gamma, \, \alpha_{\delta} \text{ is admissible relative to } \{\alpha_{\beta} : \beta < \delta\} \\ \exists r \text{ s.t. } \alpha_{\beta} \text{ is the } \beta \text{-th } r \text{-admissible} \end{array}$ 

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Use of infinite time Turing machines in generalized lemma

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 $rac{1}{2}eta +$  1-th in the sequence should not be decodable from r in  $L_{ au_{
m B}}$ 

Compatibility hypothesis: having each member admissible relative to the initial segment

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